



# Digital Logic

W. Bolton, "Mechatronics --- Electronic control systems in mechanical and electrical engineering," 5<sup>th</sup> edition, Pearson Education Limited 2012, Chap 5  
J. Edward Carryer, R. Matthew Ohline, Thomas W. Kenny, "Introduction to Mechatronic Design," Prentice Hall 2011, Chap 18  
線上學習網站 : <https://www.electronics-tutorials.ws>  
PowerPoint 中部分圖片擷取和修改自教科書和網路圖片

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機電系統原理與實驗一 ME5126 林沛群

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## Definitions

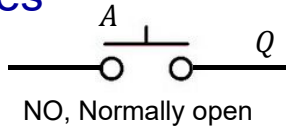
- Digital logic
  - ◆ The manipulation of binary values that uses circuits and logic gates to construct the implementation of computer operations
- Logic gates
  - ◆ The building blocks for digital electronic circuits
  - ◆ Five main components
    - AND
    - OR
    - XOR
    - NAND
    - NOR

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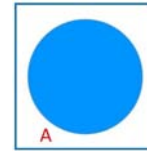
2

# Logic Gates -1

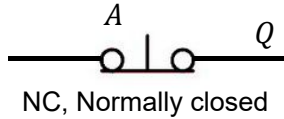
## Yes



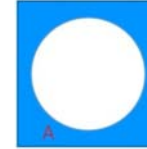
Boolean equation:  $Q = A$



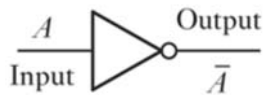
## NOT



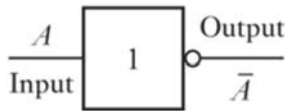
Boolean equation:  $Q = \bar{A}$



Represented by switches



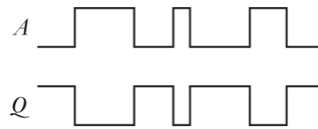
US symbols



International standard (IEEE/ANSI)

Input	Output
A	$Q = \bar{A}$
0	1
1	0

Truth table



Venn diagram:

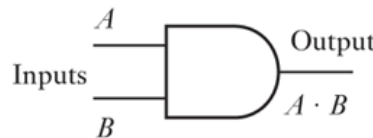
A diagram that shows all possible logical relations between a finite collection of different sets

# Logic Gates -2

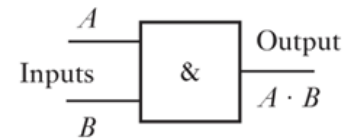
## AND



Represented by switches  
"in series"

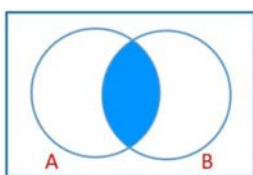


US symbols

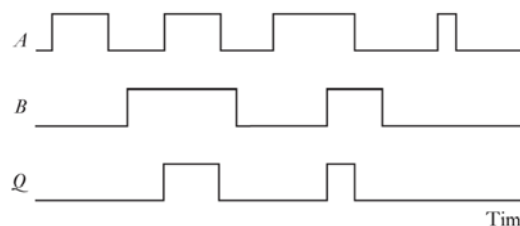


International standard  
(IEEE/ANSI)

Boolean equation:  $Q = A \cdot B$



Venn diagram

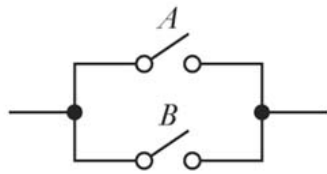


Input		Output
A	B	$Q = A \cdot B$
0	0	0
0	1	0
1	0	0
1	1	1

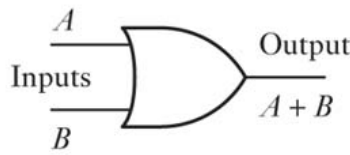
Truth table

# Logic Gates -3

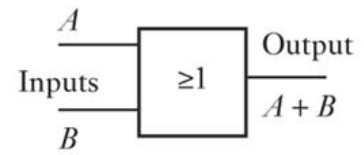
## OR



Represented by switches  
"in parallel"

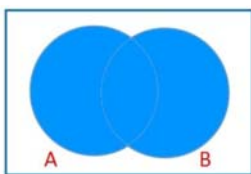


US symbols

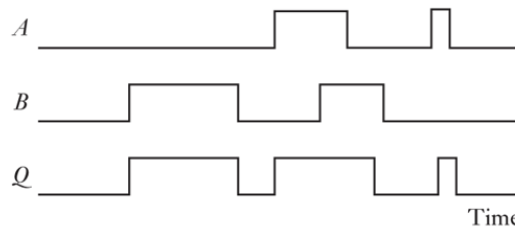


International standard  
(IEEE/ANSI)

Boolean equation:  $Q = A + B$



Venn diagram



Input		Output
A	B	$Q = A + B$
0	0	0
0	1	1
1	0	1
1	1	1

Truth table

# Boolean Algebra -1

## Operations: 0, 1, $\cdot$ (AND), $+$ (OR)

$$0 \cdot 0 = 0$$

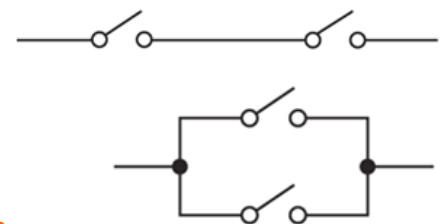
$$0 + 0 = 0$$

$$1 \cdot 1 = 1$$

$$1 + 1 = 1 \quad \leftarrow \text{Different from arithmetic operations}$$

$$1 \cdot 0 = 0 \cdot 1 = 0$$

$$1 + 0 = 0 + 1 = 1$$



## Commutative property

$$X \cdot Y = Y \cdot X$$

$$X + Y = Y + X$$

When  $Y = 0$

$$X \cdot 0 = 0$$

$$X + 0 = X$$

When  $Y = X$

$$X \cdot X = X$$

$$X + X = X$$

When  $Y = 1$

$$X \cdot 1 = X$$

$$X + 1 = 1$$

When  $Y = \bar{X}$

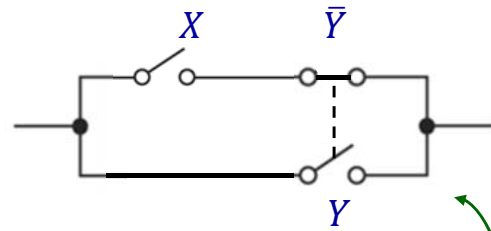
$$X \cdot \bar{X} = 0$$

$$X + \bar{X} = 1$$

## Boolean Algebra -2

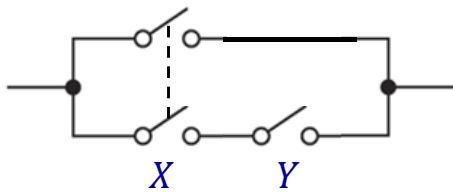
### □ Sequence of operations

- ◆ Negation (inversion)
- ◆ Multiplication
- ◆ Addition



### □ Distribution property

$$X \cdot (Y + Z) = X \cdot Y + X \cdot Z$$



### Absorption laws

$$X + X \cdot Y = X(1 + Y) = X$$

$$X \cdot (X + Y) = X \cdot X + X \cdot Y = X$$

$$(X + \bar{Y}) \cdot Y = XY$$

$$X \cdot \bar{Y} + Y = X + Y$$

$$X \cdot Y + \bar{Y} = X + \bar{Y}$$

## Boolean Algebra -3

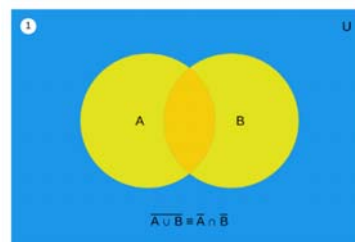
### □ De Morgan's Laws

$$\overline{(X + Y + Z + \dots)} = \bar{X} \cdot \bar{Y} \cdot \bar{Z} \cdot \dots$$

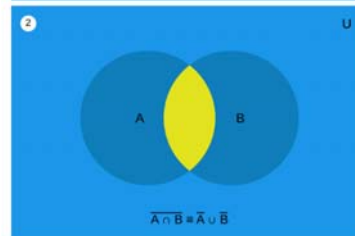
$$\overline{(X \cdot Y \cdot Z \cdot \dots)} = \bar{X} + \bar{Y} + \bar{Z} + \dots$$

- ◆ Two variables

$$\overline{(A + B)} = \bar{A} \cdot \bar{B}$$



$$\overline{(A \cdot B)} = \bar{A} + \bar{B}$$



## Boolean Algebra -4

### □ Multiple variables

" · " is omitted

$$X + Y + Z = (X + Y) + Z = X + (Y + Z)$$

$$XYZ = (XY)Z = X(YZ)$$

$$(X + Y)(X + Z) = XX + XY + XZ + YZ = X(1 + Y + Z) + YZ = X + YZ$$

$$(X + Y)(X + Z)(Z + \bar{X}) = (X + Y)(Z + \bar{X})$$

$$(X + Y)(Z + \bar{X}) = XZ + \bar{X}Y$$

$$XY + YZ + Z\bar{X} = XY + Z\bar{X}$$



How to prove? Use this as the example

## Boolean Algebra -5

### □ Algebraic proof

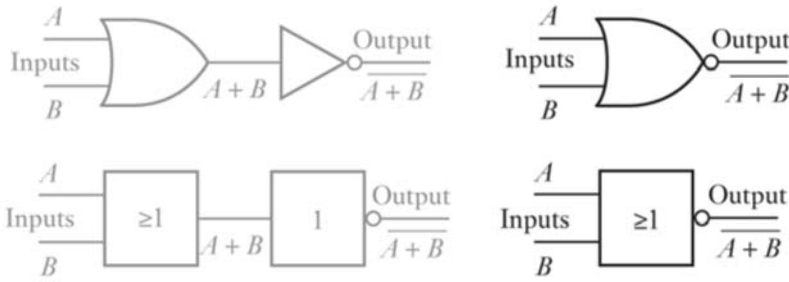
$$\begin{aligned} XY + YZ + Z\bar{X} &= XY + 1YZ + Z\bar{X} = XY + (X + \bar{X})YZ + Z\bar{X} \\ &= XY + XYZ + \bar{X}YZ + Z\bar{X} = XY(1 + Z) + (Y + 1)Z\bar{X} \\ &= XY + Z\bar{X} \end{aligned}$$

### □ Brute force

X	Y	Z	XY	YZ	Z $\bar{X}$	XY + YZ + Z $\bar{X}$	XY + Z $\bar{X}$
0	0	0	0	0	0	0	0
0	0	1	0	0	1	1	1
0	1	0	0	0	0	0	0
0	1	1	0	1	1	1	1
1	0	0	0	0	0	0	0
1	0	1	0	0	0	0	0
1	1	0	1	0	0	1	1
1	1	1	1	1	0	1	1

# Logic Gates -4

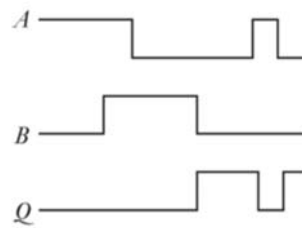
□ NOR = OR + NOT



US symbols

International standard (IEEE/ANSI)

Boolean equation:  $Q = \overline{A+B} = \bar{A} \cdot \bar{B}$



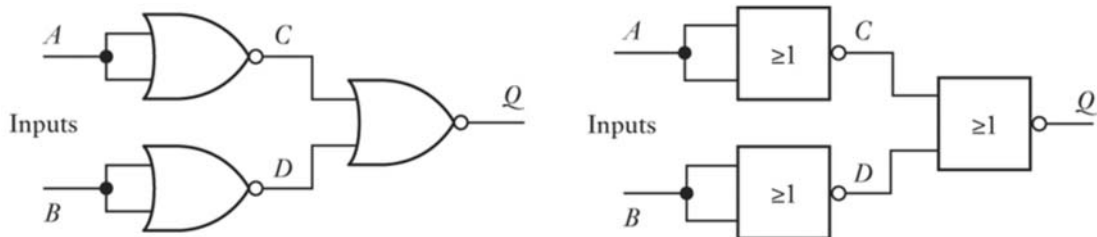
Input		Output
A	B	$Q = \overline{A+B}$
0	0	1
0	1	0
1	0	0
1	1	0

Truth table

# Logic Gates -5

□ NOR: a universal gate

◆ AND: Using three NOR gates



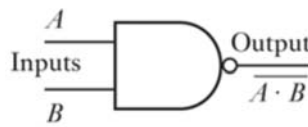
Input				Output
A	B	C	D	$Q = A \cdot B$
0	0	1	1	0
0	1	1	0	0
1	0	0	1	0
1	1	0	0	1

$$\begin{aligned}
 Q &= \overline{\overline{C+D}} \\
 &= \overline{\bar{C} \cdot \bar{D}} \\
 &= \overline{\bar{A} + \bar{A} \cdot \bar{B} + \bar{B}} \\
 &= A \cdot B \\
 &= \text{AND}
 \end{aligned}$$

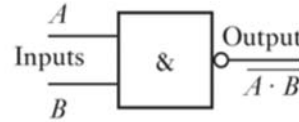
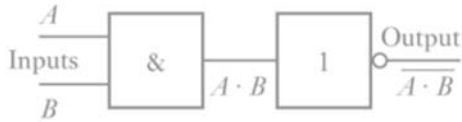
Using NOR to make other gates such as NOT, OR, NAND, XOR...?

## Logic Gates -6

- NAND = AND + NOT

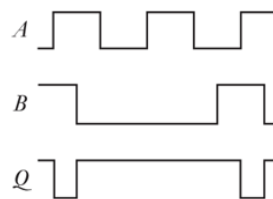


US symbols



International standard  
(IEEE/ANSI)

Boolean equation:  $Q = \overline{A \cdot B} = \bar{A} + \bar{B}$

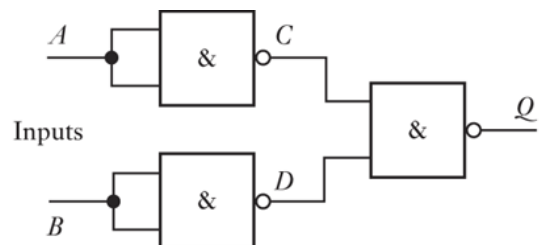
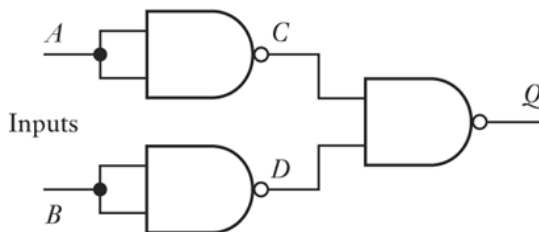


Input		Output
A	B	$Q = \overline{A \cdot B}$
0	0	1
0	1	1
1	0	1
1	1	0

Truth table

## Logic Gates -7

- NAND: a universal gate
  - ◆ Duality: NOR & NAND
  - ◆ OR: Using three NAND gates

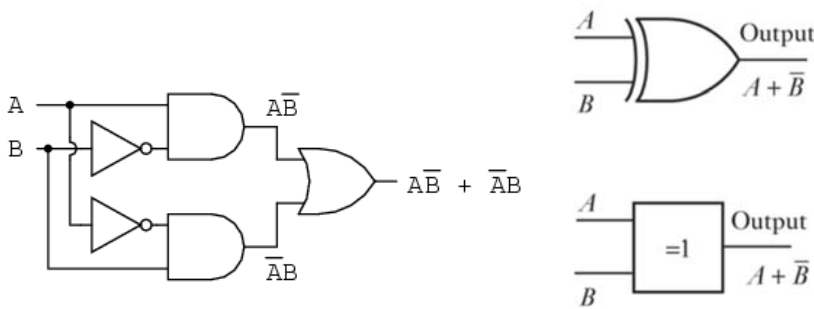


Input				Output
A	B	C	D	$Q = A + B$
0	0	1	1	0
0	1	1	0	1
1	0	0	1	1
1	1	0	0	1

$$\begin{aligned}
 Q &= \overline{\overline{C \cdot D}} \\
 &= \overline{\overline{C} + \overline{D}} \\
 &= \overline{\overline{A \cdot A} + \overline{B \cdot B}} \\
 &= A + B \\
 &= \text{OR}
 \end{aligned}$$

# Logic Gates -8

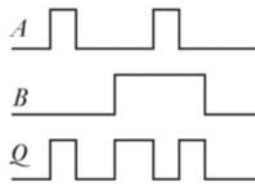
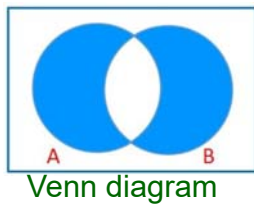
## XOR



US symbols

International standard (IEEE/ANSI)

Boolean equation:  $Q = A \oplus B = A \cdot \bar{B} + \bar{A} \cdot B$



Input		Output
A	B	$Q = A \oplus B$
0	0	0
0	1	1
1	0	1
1	1	0

Truth table

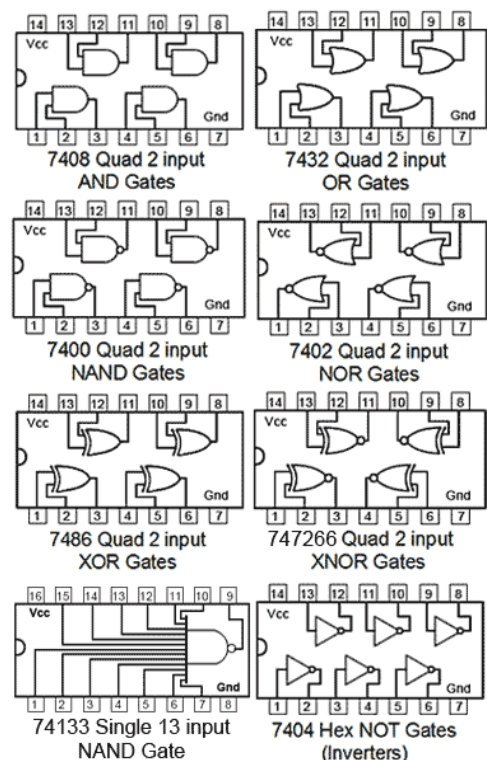
# Logic Families and Integrated Circuits -1

## Families

- ◆ Transistor-transistor logic (TTL)
- ◆ Complementary metaloxide semiconductor (CMOS)
- ◆ Emitter-coupled logic (ECL)

## Integrated circuits

- ◆ 74XX series (TTL)
  - 74YYXX, YY=(none), LS, F...
- ◆ 4000 series (CMOS)
- ◆ 74CXX, 74HCXX (CMOS)





## Logic Families and Integrated Circuits -2

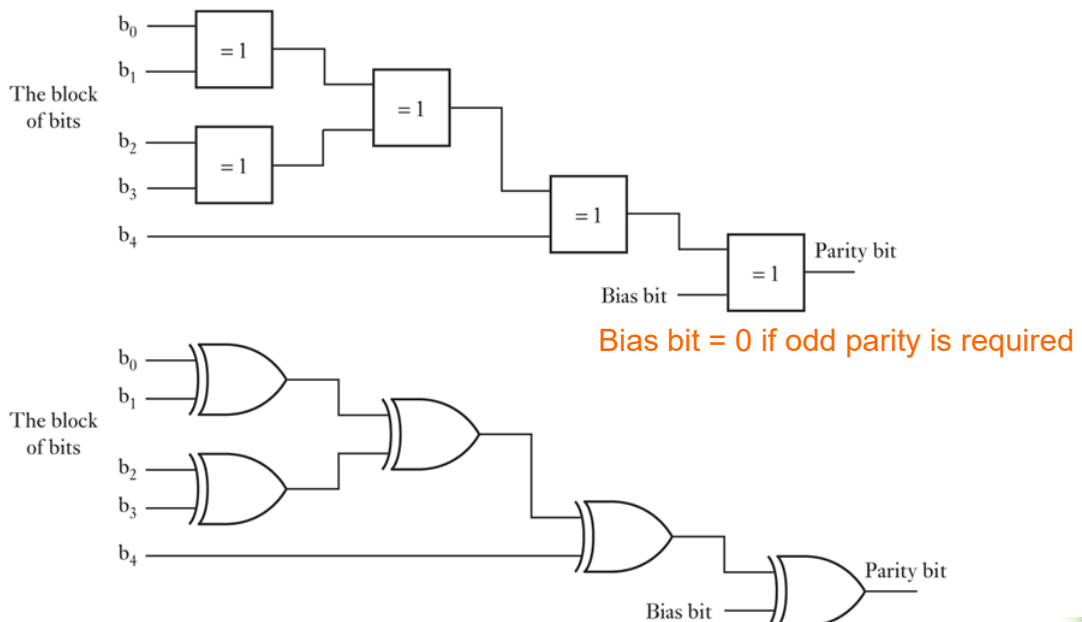
### □ Parameters

- ◆ Logic level
- ◆ Noise immunity
- ◆ Fan-out: The number of gate inputs that can be driven by a standard gate output while maintaining the desired LOW or HIGH levels
- ◆ Current-sourcing or current-sinking action
- ◆ Propagation delay time
- ◆ Power consumption

## Applications of Logic Gates -1

### □ Parity generators

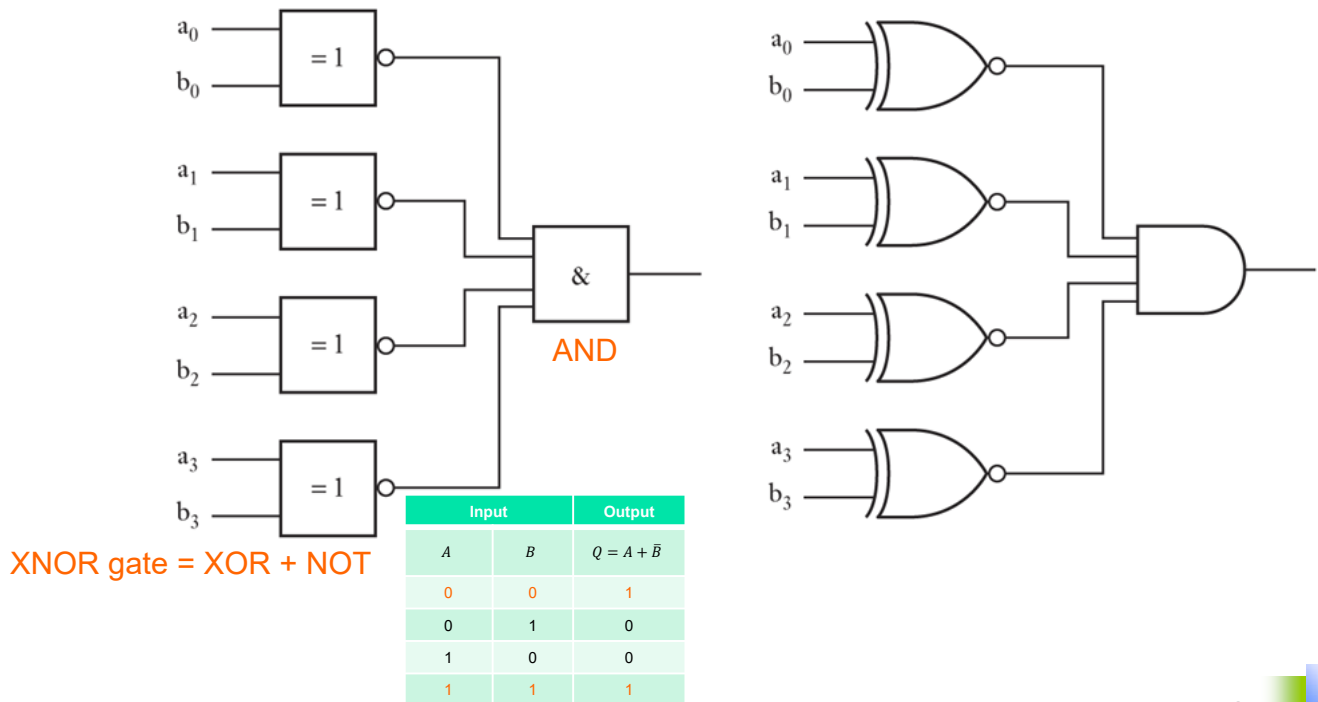
- ◆ Using XOR: output 0 if all the inputs are 0 or all are 1; otherwise output 1



## Applications of Logic Gates -2

### □ Digital comparator

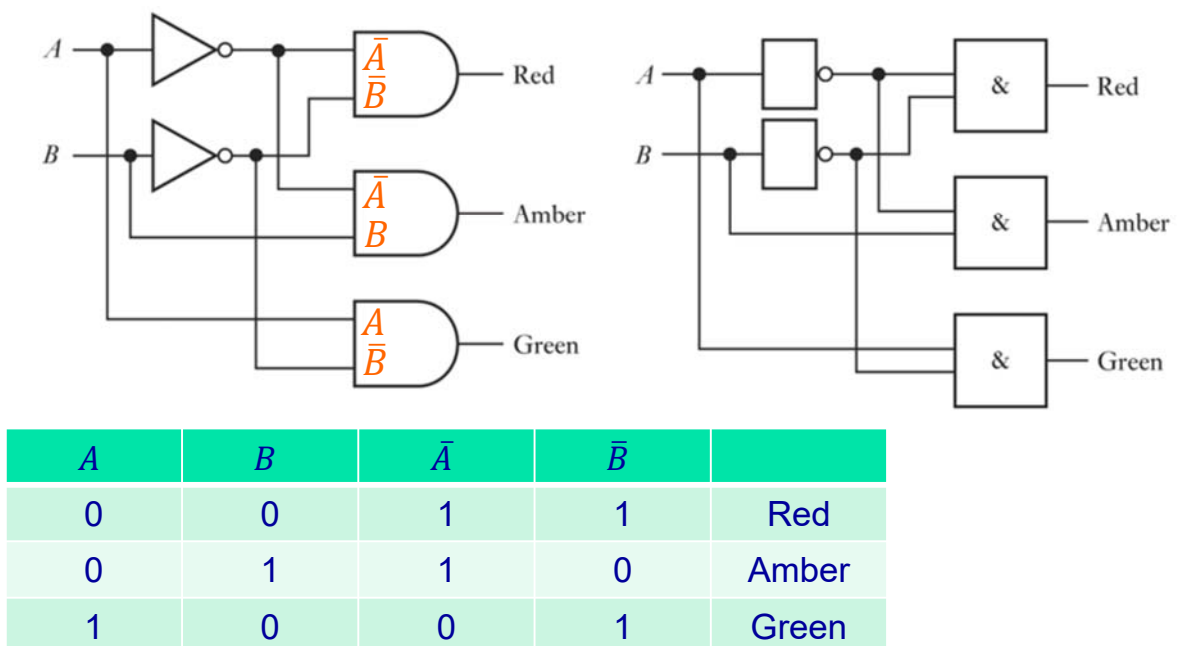
- ◆ Comparing two digital words to determine if they are exactly equal



## Applications of Logic Gates -3

### □ Coder / decoder

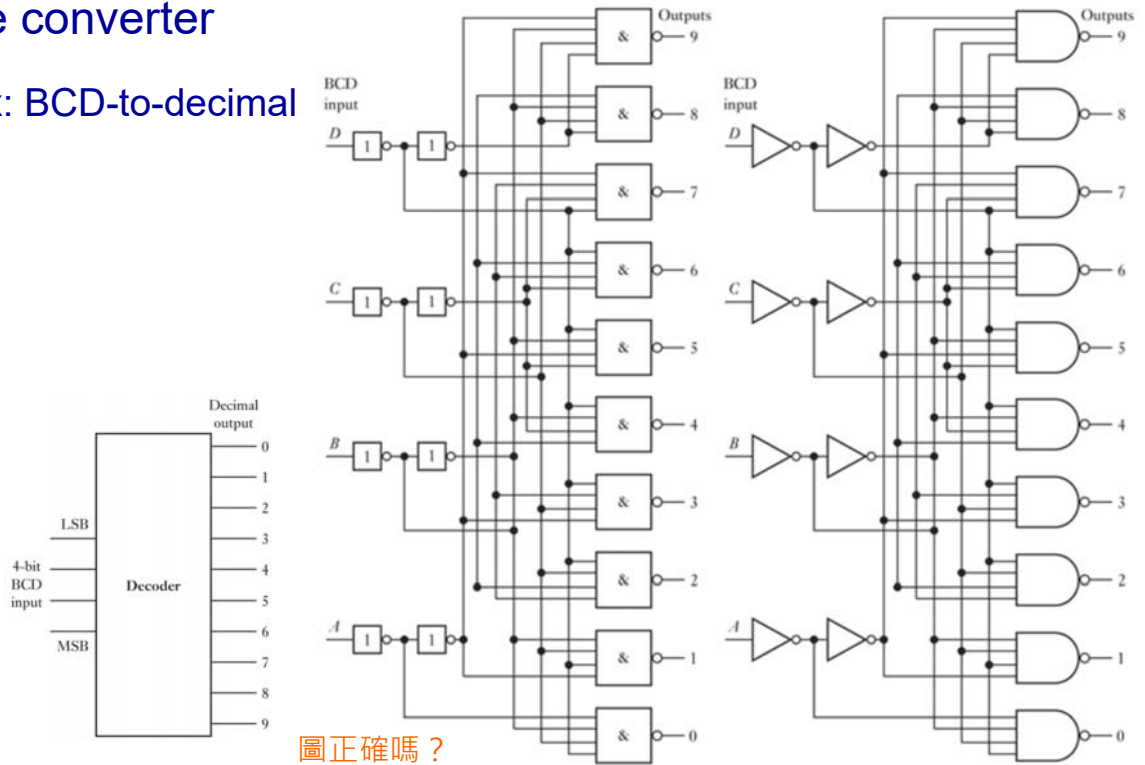
- ◆ Ex: Traffic lights



# Applications of Logic Gates -4

## Code converter

### Ex: BCD-to-decimal



### Ex: 74LS244 --- 4-bit BCD to seven segments of a display

# Applications of Logic Gates -5

## Number systems

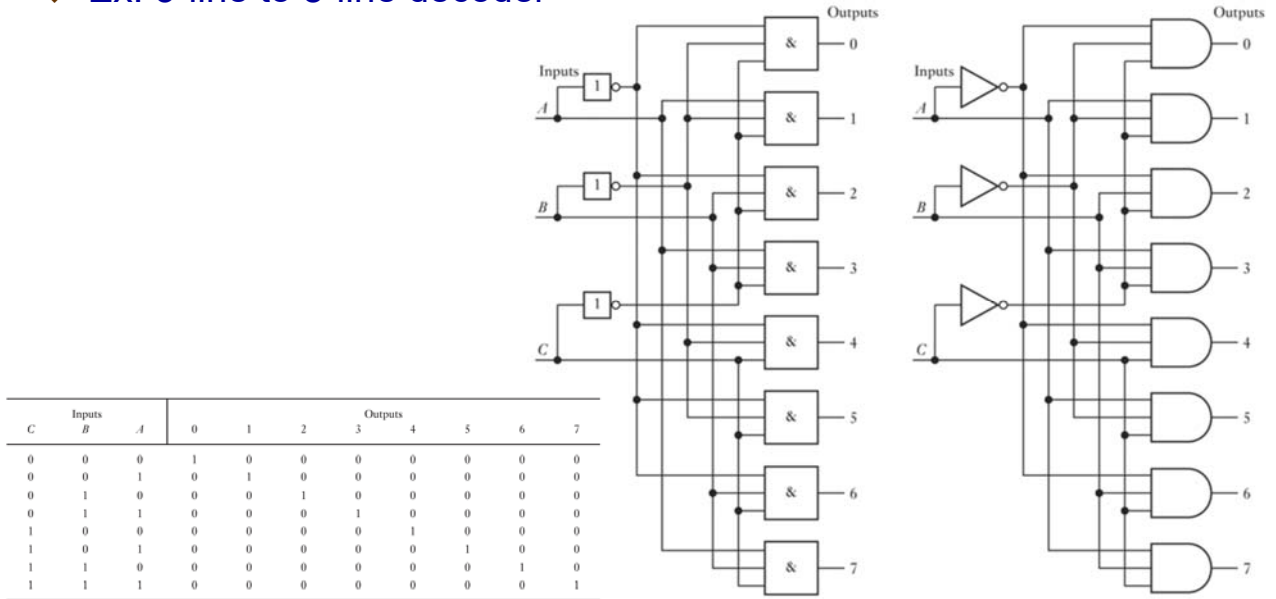
Binary-coded decimal system

Decimal	Binary	BCD	Octal	Hexadecimal
0	0000	0000 0000	0	0
1	0001	0000 0001	1	1
2	0010	0000 0010	2	2
3	0011	0000 0011	3	3
4	0100	0000 0100	4	4
5	0101	0000 0101	5	5
6	0110	0000 0110	6	6
7	0111	0000 0111	7	7
8	1000	0000 1000	10	8
9	1001	0000 1001	11	9
10	1010	0001 0000	12	A
11	1011	0001 0001	13	B
12	1100	0001 0010	14	C
13	1101	0001 0011	15	D
14	1110	0001 0100	16	E
15	1111	0001 0101	17	F

# Applications of Logic Gates -6

## Decoder

### Ex: 3-line to 8-line decoder

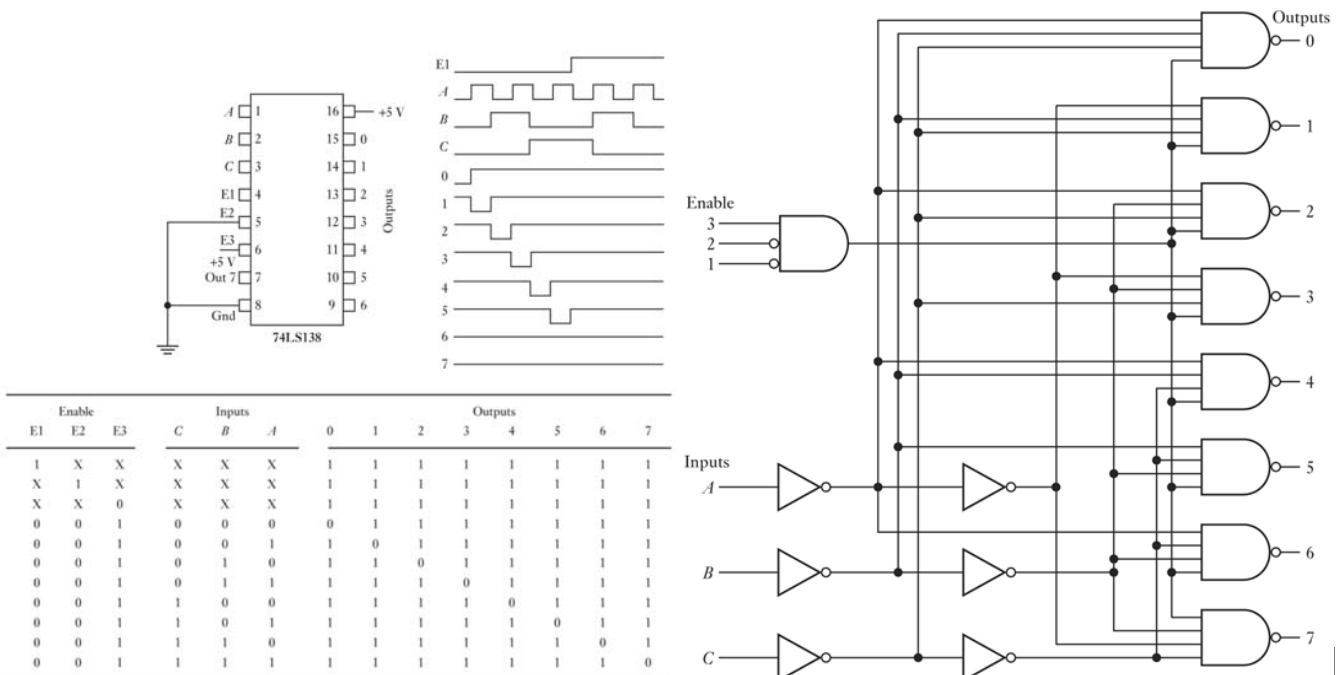


### Ex: 74LS138

# Applications of Logic Gates -7

## Decoder

### Ex: 74LS138, 3-line to 8-line decoder with ENABLE inputs



## Simplification of Binary Functions -1

- Where do these complicate functions come from?
  - ◆ The result of translating the practical requirements of a given control problem into the logic language of Boolean algorithm
  
- Why do we want to simplify?
  - ◆ Save number of logic gates
  - ◆ Easier to understand the consequence

## Simplification of Binary Functions -2

- Method 1 – Algebraic method
- Ex:

$$\begin{aligned} & ABCD + ABC\bar{C} + AB\bar{D} && 2 \text{ NOT, } 3 \text{ AND, } 1 \text{ OR} \\ & = AB(CD + \bar{C} + \bar{D}) \\ & = AB(\bar{C} + D + \bar{D}) \\ & = AB(\bar{C} + 1) \\ & = AB1 \\ & = AB && 1 \text{ AND} \end{aligned}$$

## Simplification of Binary Functions -3

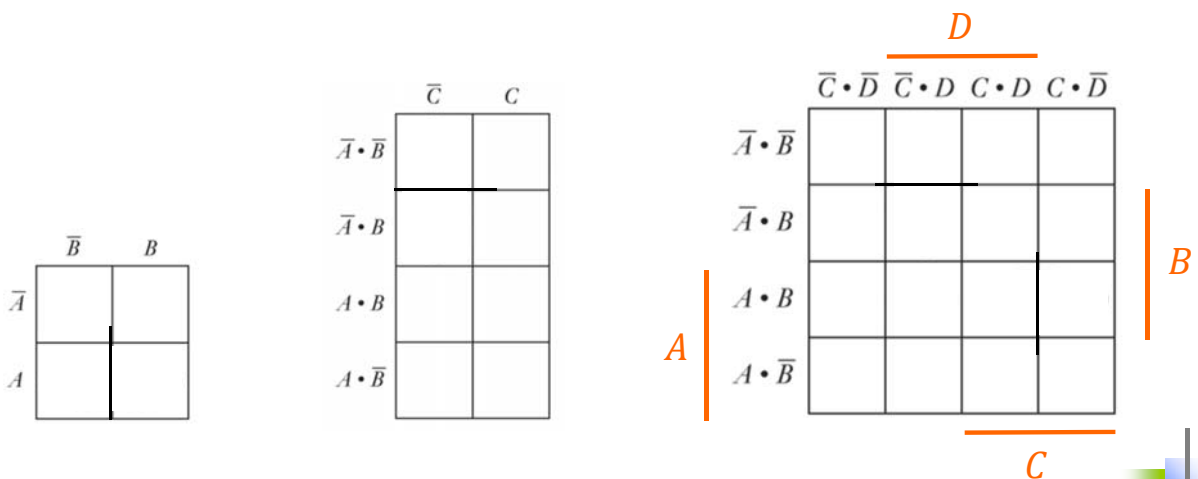
### □ Method 2 – Karnaugh Map

- ◆ A graphical method that can be used to produce simplified Boolean expressions from sums of products obtained from truth tables
- ◆ Advantages
  - No need tricks
  - Faster and simpler
  - Always produces the simplest possible result
  - Include “don’t care” conditions

## Karnaugh Map -1

### □ Generation of map

- Any two adjacent squares MUST differ in only one variables (top-bottom left-right edges also considered adjacent)



## Karnaugh Map -2

### □ Procedure

- ◆ Fill "1" into the map, then "0" for the rest
- ◆ Loop adjacent "1" into cells of 2,4,8 cells
  - Each "1" must be covered by at least one cell
  - Cell number: as small as possible
  - Cell size: as large as possible

	$\bar{B}$	$B$
$\bar{A}$		
$A$	1	1

$$T = AB + A\bar{B}$$

	$\bar{C}$	$C$
$\bar{A}\cdot\bar{B}$	1	
$\bar{A}\cdot B$	1	
$A\cdot B$		
$A\cdot\bar{B}$		

$$T = \bar{A}\bar{B}\bar{C} + \bar{A}B\bar{C}$$

$$T = ABCD + ABC\bar{D} + \bar{A}B\bar{C}D + \bar{A}\bar{B}\bar{C}D$$

	$D$			
	$\bar{C}\cdot\bar{D}$	$\bar{C}\cdot D$	$C\cdot D$	$C\cdot\bar{D}$
$\bar{A}\cdot\bar{B}$		1		
$\bar{A}\cdot B$		1		
$A\cdot B$			1	1
$A\cdot\bar{B}$				

## Karnaugh Map -3

- ◆ Write new logic

$$T = AB + A\bar{B} = A$$

	$\bar{B}$	$B$
$\bar{A}$		
$A$	1	1

$$T = \bar{A}\bar{B}\bar{C} + \bar{A}B\bar{C} = \bar{A}\bar{C}$$

	$\bar{C}$	$C$
$\bar{A}\cdot\bar{B}$	1	
$\bar{A}\cdot B$	1	
$A\cdot B$		
$A\cdot\bar{B}$		

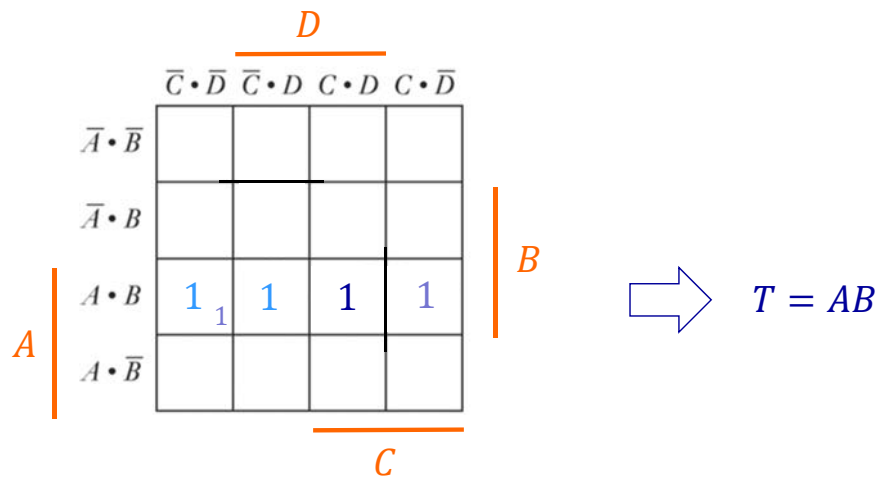
$$T = ABCD + ABC\bar{D} + \bar{A}B\bar{C}D + \bar{A}\bar{B}\bar{C}D = \bar{A}\bar{C}D + ABC$$

	$D$			
	$\bar{C}\cdot\bar{D}$	$\bar{C}\cdot D$	$C\cdot D$	$C\cdot\bar{D}$
$\bar{A}\cdot\bar{B}$		1		
$\bar{A}\cdot B$		1		
$A\cdot B$			1	1
$A\cdot\bar{B}$				

# Karnaugh Map -4

- Ex: Revisit the same example

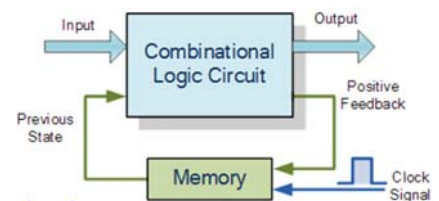
$$ABCD + ABC\bar{C} + AB\bar{D}$$



# Sequential Logic -1

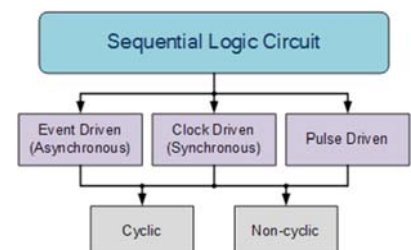
- A type of logic circuit whose output depends on

- The present value of its input signals
- The sequence of past inputs, “memory”
- Generally a bistable device which can have their output or outputs set in one of two basic states



- Classification

- Event driven, asynchronous circuits that change state immediately when enabled
- Clock driven, synchronized to a specific clock signal
- Pulse driven, a combination of the two



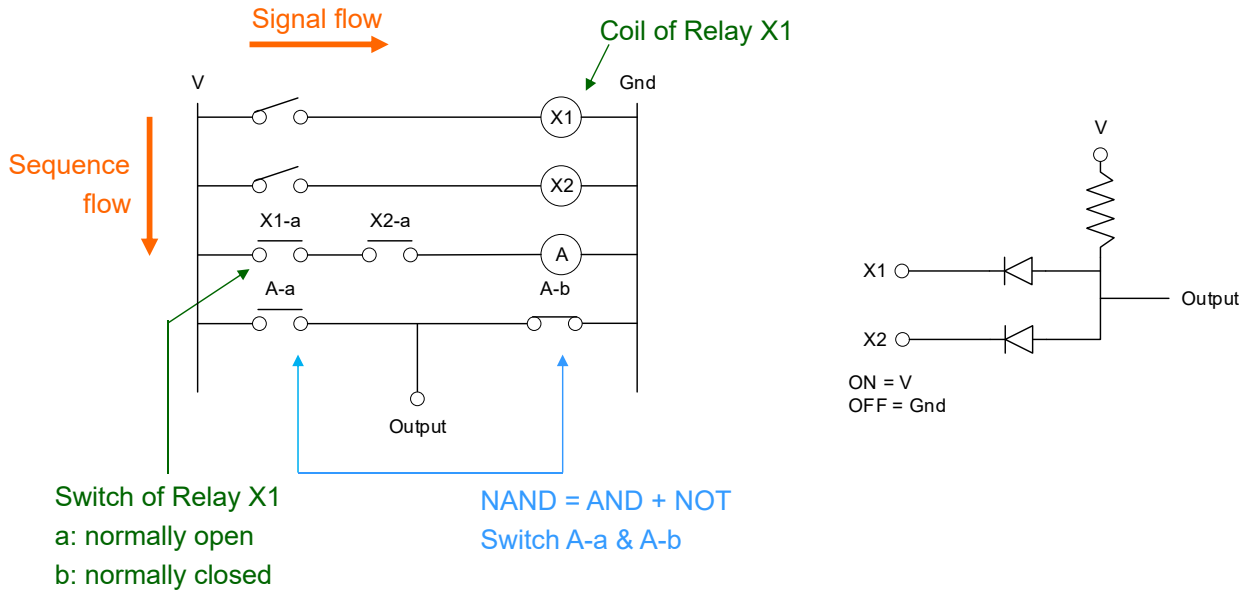


# Sequential Logic -2

## Basic gates using relays

### AND

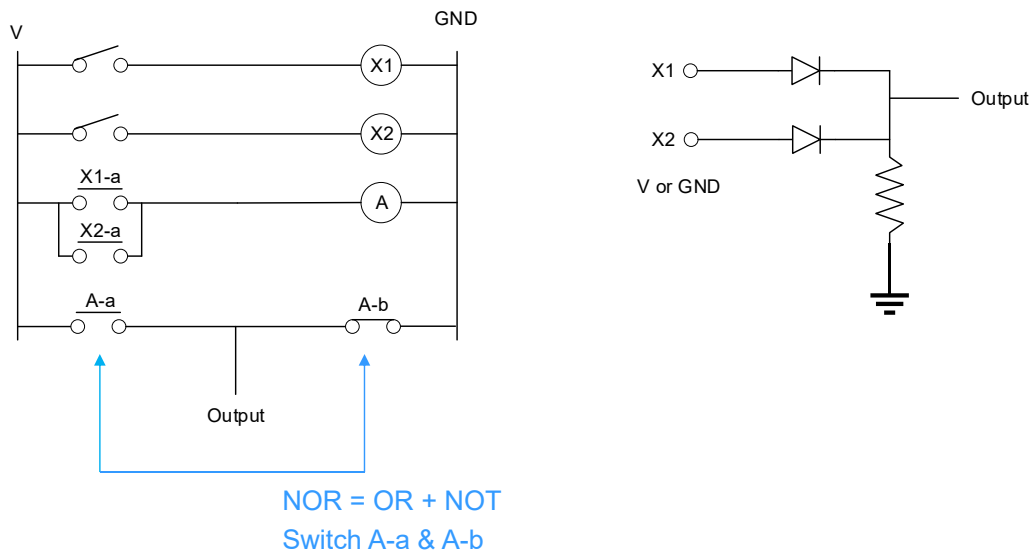
Ladder diagram



# Sequential Logic -3

## Basic gates using relays

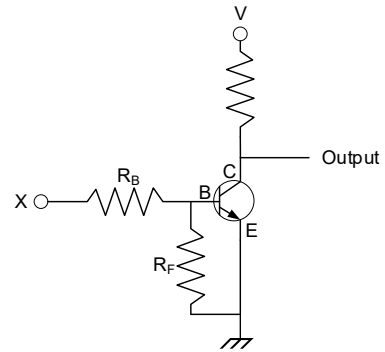
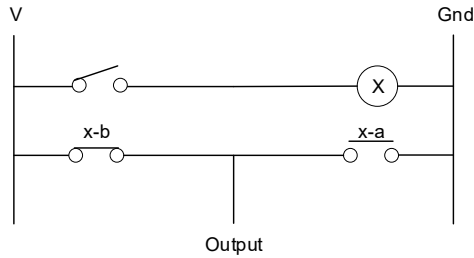
### OR



# Sequential Logic -4

## Basic gates using relays

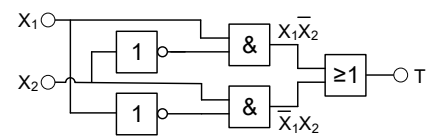
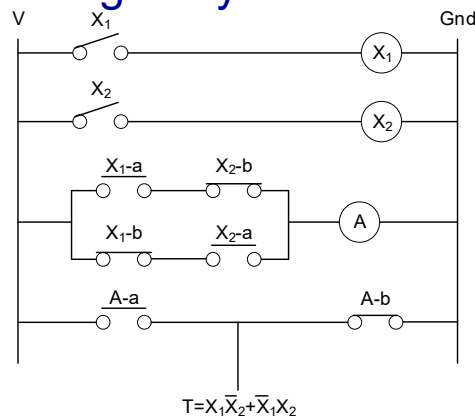
### NOT



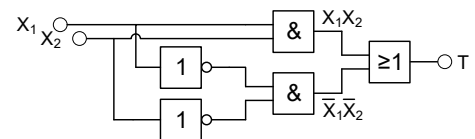
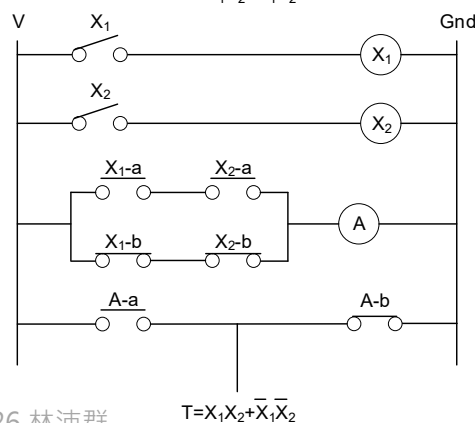
# Sequential Logic -5

## Basic gates using relays

### XOR

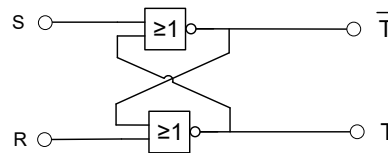
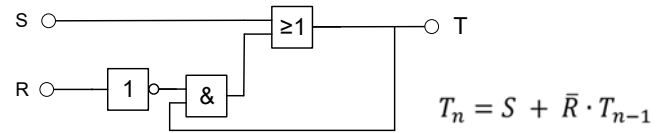
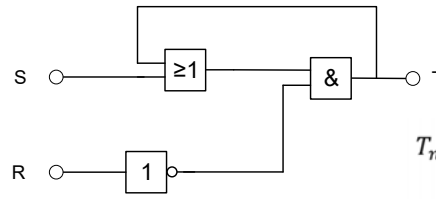
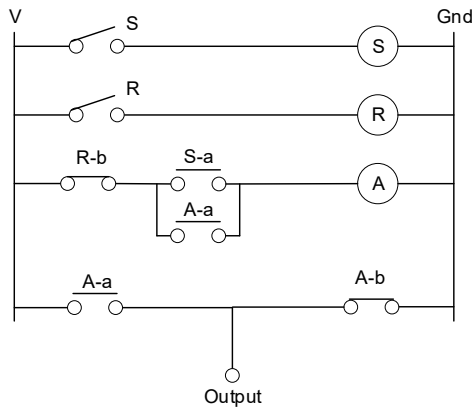


### 一致電路



# Sequential Logic -6

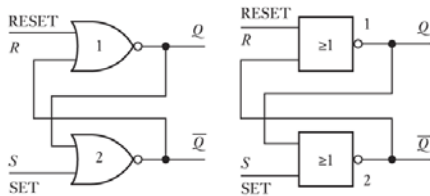
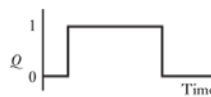
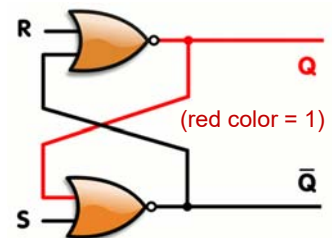
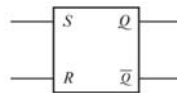
- Latching function 自我保持, a SR flip-flop circuit



# Sequential Logic -7

- SR (set-reset) flip-flop

$$Q_{n+1} = \bar{R} \cdot (Q_n + S)$$



$Q_n$	$S$	$R$	$Q_{n+1}$	Action
0	0	--	0	Hold state
0	1	0	1	Set
1	0	1	0	Reset
1	--	0	1	Hold state

Excitation table (-- doesn't matter)

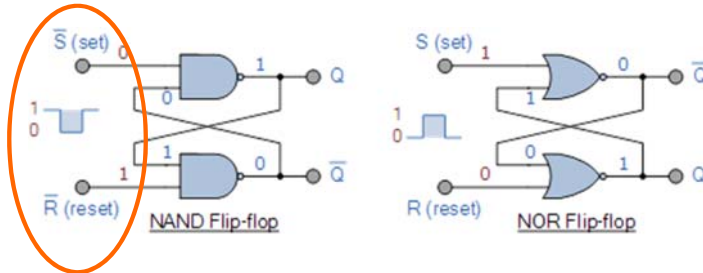
$Q_n$	$S$	$R$	$Q_{n+1}$
0	0	0	0
0	1	0	1
0	0	1	0
0	1	1	X
1	0	0	1
1	1	0	1
1	0	1	0
1	1	1	X

Truth table

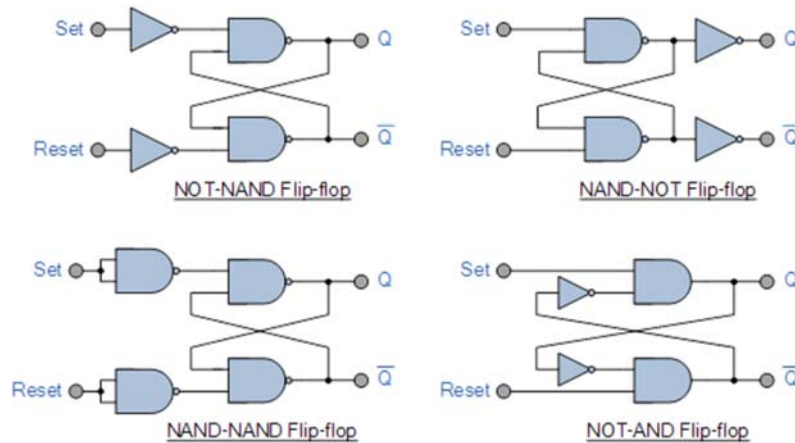
# Sequential Logic -8

## SR flip-flop

### Basic forms



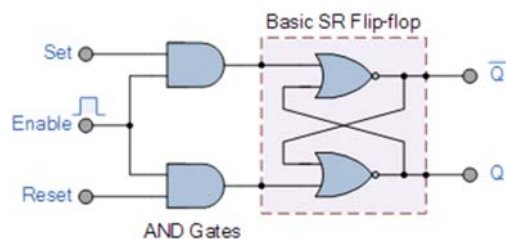
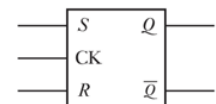
### Active HIGH



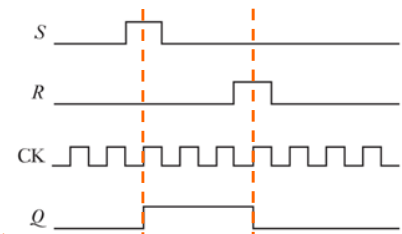
# Sequential Logic -9

## Clocked (Gated) SR flip-flop

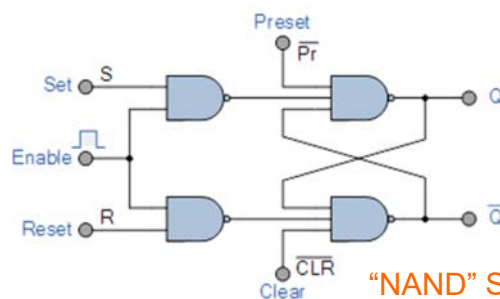
- A synchronous system, to regulate the set/reset timings



“NOR” SR – HIGH inputs



## SR Flip-flop with Preset and Clear Inputs



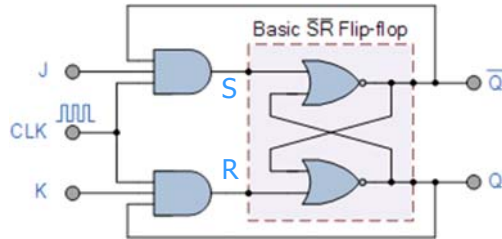
“NAND” SR – LOW inputs

# Sequential Logic -10

## □ Clocked JK flip-flop

◆ = Clocked RS + “flip” (when  $J \cdot K = 1$ )

$$Q_{n+1} = J \cdot \overline{Q_n} + \overline{K} \cdot Q_n$$



$Q_n$	$J$	$K$	$Q_{n+1}$	Action
$Q_n$	0	0	$Q_n$	Hold state
0	1	0	1	Set
1	0	1	0	Reset
$Q_n$	1	1	$\overline{Q_n}$	Toggle

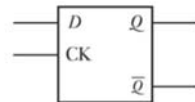
Excitation table

$Q_n$	$J$	$K$	$Q_{n+1}$
0	0	0	0
0	1	0	1
0	0	1	0
0	1	1	1
1	0	0	1
1	1	0	1
1	0	1	0
1	1	1	0

Truth table

# Sequential Logic -11

## □ D flip-flop

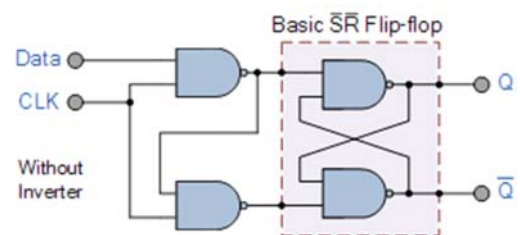
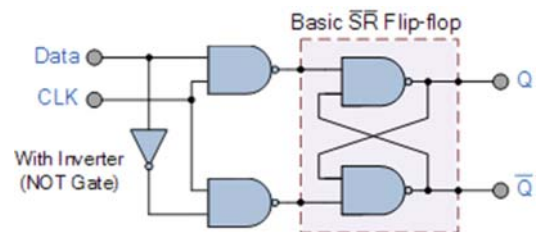
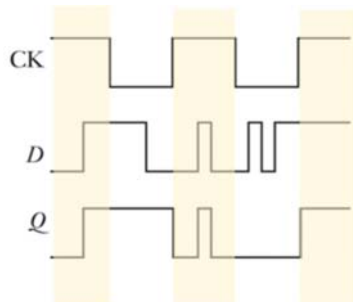
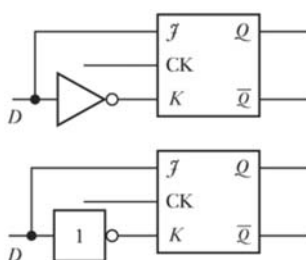


◆ = SR/JK + NOT

◆ One control input, forcing  $S \cdot R = 0$  or  $J \cdot K = 0$

◆  $D = 1$  set (S);  $D = 0$  reset (R)

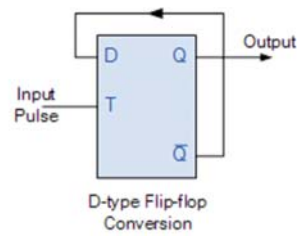
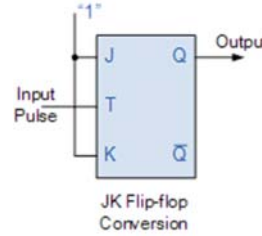
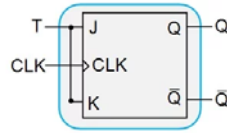
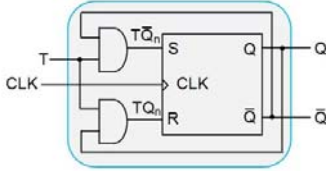
$$Q_{n+1} = D$$



# Sequential Logic -12

## T flip-flop

$$Q_{n+1} = T \oplus Q_n = T \cdot \overline{Q_n} + \overline{T} \cdot Q_n$$



$$J = K = T$$

“flip” when  $T = 1$

$$J \cdot K = 1$$

“flip” when  $T = 1$

$Q_n$	$T$	$Q_{n+1}$	Action
$Q_n$	0	$Q_n$	Hold state
$Q_n$	1	$\overline{Q_n}$	Toggle

Characteristic table

$Q_n$	$T$	$Q_{n+1}$
0	0	0
0	1	1
1	0	1
1	1	0

Truth table

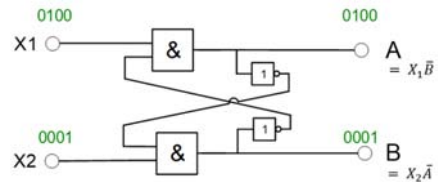
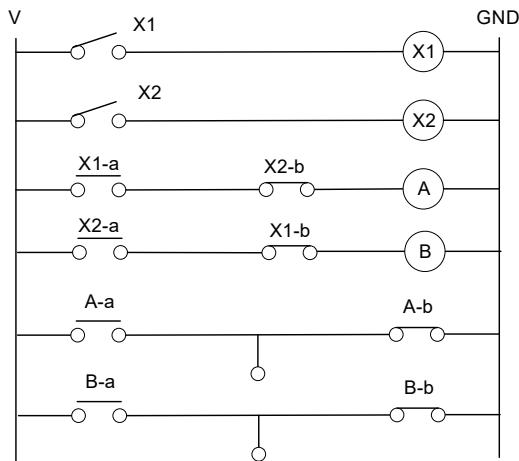
# Sequential Logic -13

## Popular flip flop ICs

Device Number	Device Description
74LS73A	Dual Negative-edge Triggered JK Flip-flop with Clear
74LS74	Dual Positive-edge Triggered D-type with Preset and Clear
74LS75	Quad D-type Bistable Latch with Enable
74LS76	Dual Pulse Triggered JK Flip-flops with Preset and Clear
74LS107	Dual JK Flip-Flop with Clear
74LS111	Dual Master-slave JK Flip-flop with Clear
74LS175	Dual Positive-edge Triggered D-type with Clear
74LS279	Quad SR Latches with Active-LOW inputs

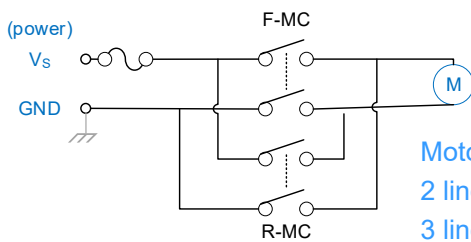
# Sequential Logic -14

## □ Interlock function 互鎖電路

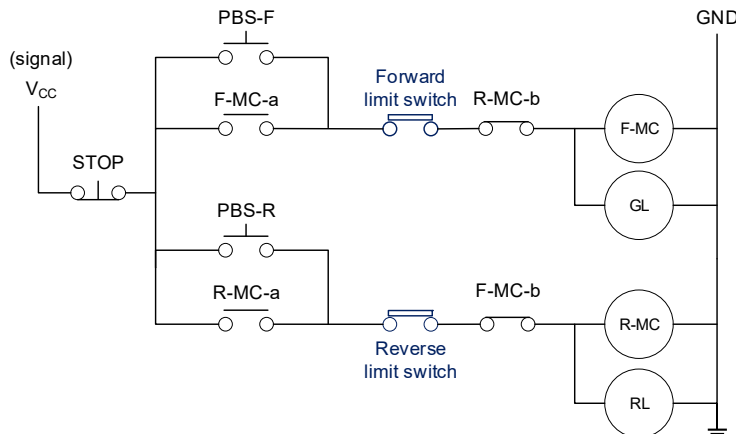
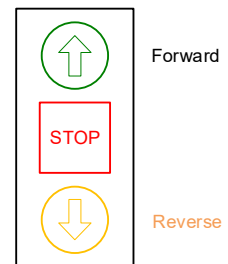


# Sequential Logic -15

## □ Motor forward/reverse control using relays



Motor:  
2 lines: DC, AC (110V – 1p, 220V – 1p)  
3 lines: AC (220V 3p)



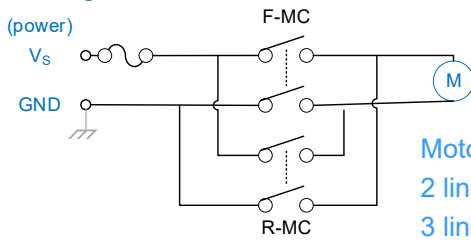
Features:  
Power-Signal separation  
Latching  
Forward/reverse interlock  
Auto-stop (limit switches)

Multiple sets of control switches?  
(onsite & remote)  
Forward/reverse – parallel  
Stop - serial

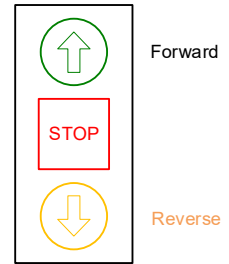
# Sequential Logic -16

## Motor forward/reverse control using logic gates

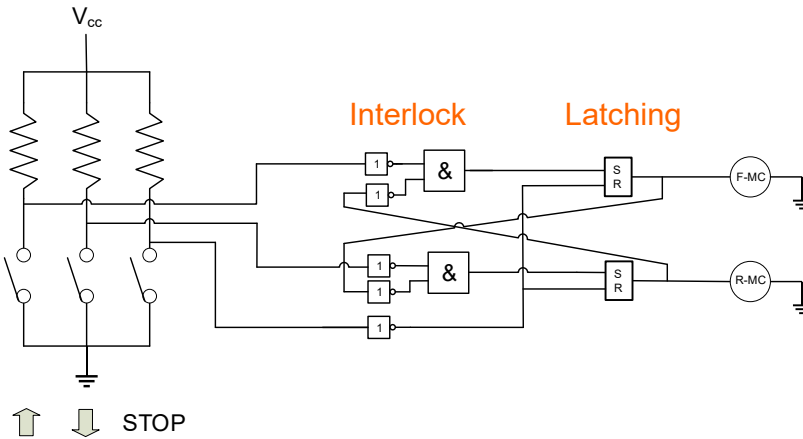
### Driving circuit



Motor:  
 2 lines: DC, AC (110V – 1p, 220V – 1p)  
 3 lines: AC (220V 3p)



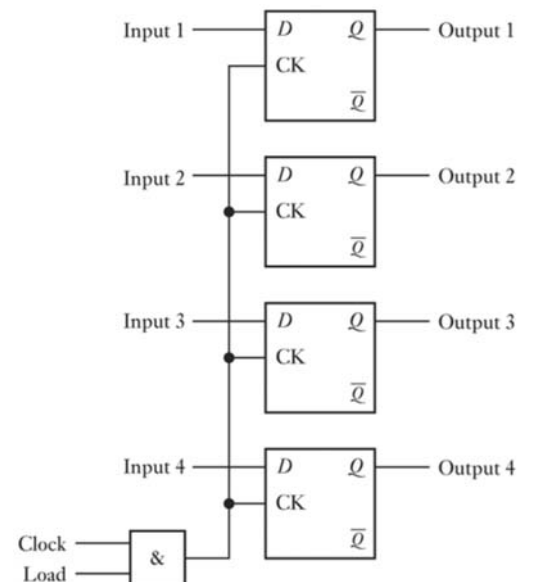
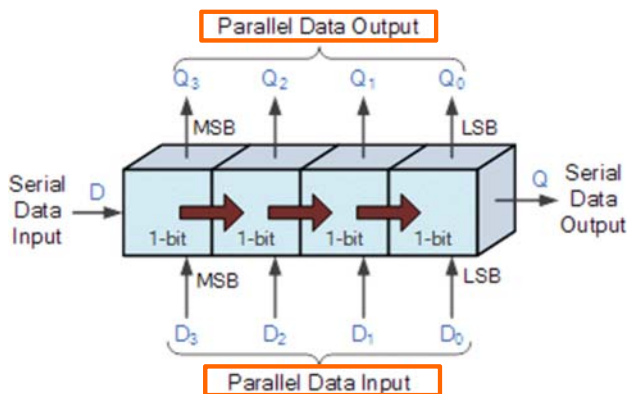
### Control circuit



# Registers -1

## Register

- ◆ A set of memory elements which is used to hold information until it is needed



Use "load" to control update timing

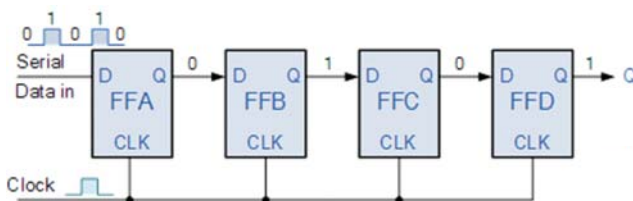


# Registers -2

## □ Shift register

- ◆ A **cascade** of flip flops, sharing the same clock, in which the output of each flip-flop is connected to the "data" input of the next flip-flop in the chain, resulting in a circuit that shifts by one position the "bit array" stored in it

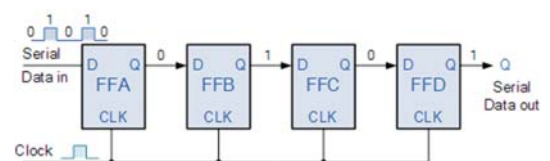
Clock Pulse No	QA	QB	QC	QD
0	0	0	0	0
1	1	0	0	0
2	0	1	0	0
3	0	0	1	0
4	0	0	0	1
5	0	0	0	0



# Registers -3

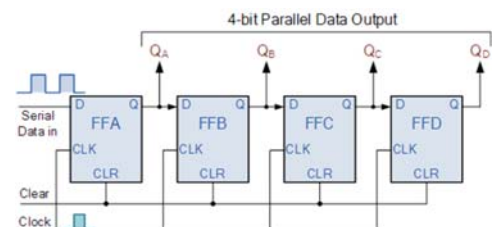
## □ Configurations

- ◆ Serial-in serial-out (SISO)

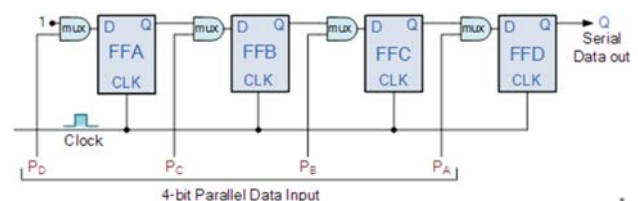


- ◆ Serial-in parallel-out (SIPO)

- Conversion between serial and parallel interfaces



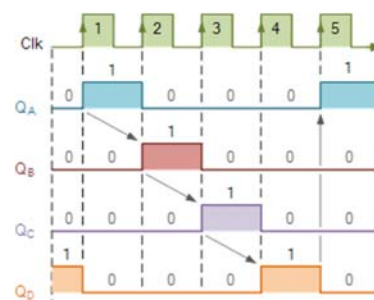
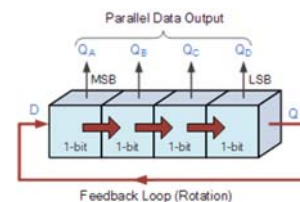
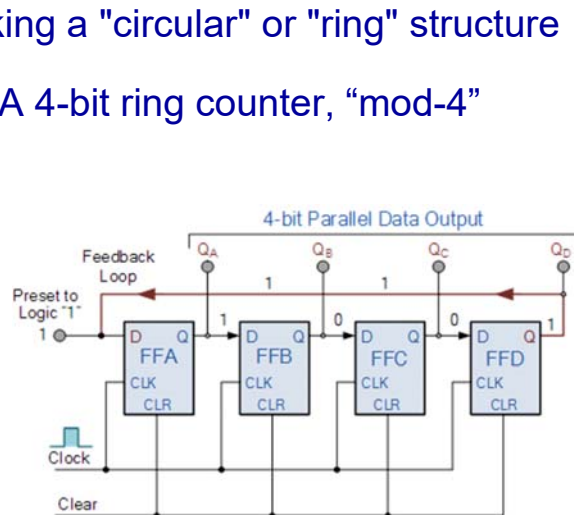
- ◆ Parallel-in serial-out (PISO)



# Ring Counter -1

## □ Ring counter

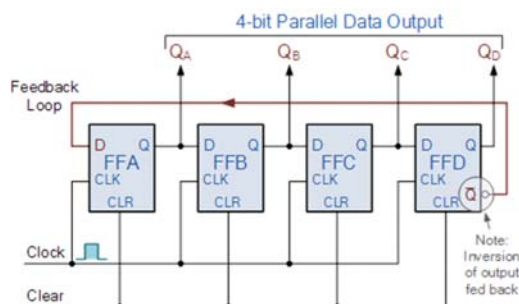
- ◆ A type of counter composed of flip-flops connected into a shift register, with the output of the last flip-flop fed to the input of the first, making a "circular" or "ring" structure
- ◆ Ex: A 4-bit ring counter, "mod-4"



# Ring Counter -2

## □ Johnson (Twisted) ring counter

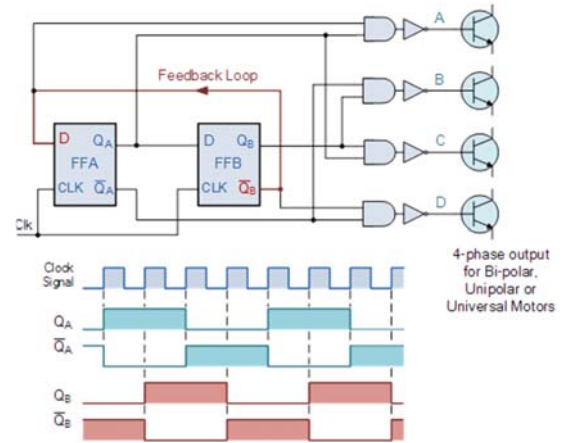
- ◆ Connecting  $\bar{Q}$  of the last flip-flop back to the input  $D$  of the first flip-flop
- ◆  $n$ -stage counter  $\rightarrow 2n$  states
- ◆ Ex: A 4-bit Johnson ring counter, "mod-8"
  - Four "0"s and four "1"s, 8-bit pattern



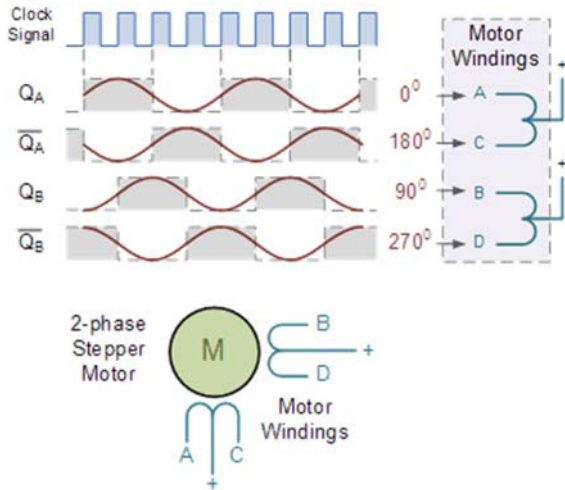
Clock Pulse No	FFA	FFB	FFC	FFD
0	0	0	0	0
1	1	0	0	0
2	1	1	0	0
3	1	1	1	0
4	1	1	1	1
5	0	1	1	1
6	0	0	1	1
7	0	0	0	1

# Ring Counter -3

- 2-bit quadrature generator



- Stepper motor control



Output	A	B	C	D
$Q_A + \bar{Q}_B$	1	0	0	0
$\bar{Q}_A + Q_B$	0	1	0	0
$Q_A + Q_B$	0	0	1	0
$\bar{Q}_A + \bar{Q}_B$	0	0	0	1

2-bit Quadrature Oscillator, Count Sequence

# The End

- Questions?

